

Reasoning about privacy properties of biometric system architectures in the presence of information leakage

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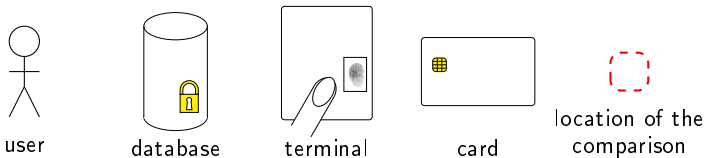
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Analysis of Biometric Systems

Analysis of Biometric Systems

Introducing biometric systems

▶ typical components



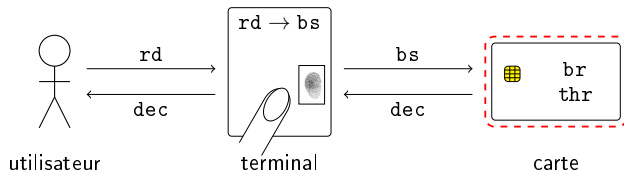
▶ typical data

- ▶ **br**: biometric reference (of an enrolled user)
- ▶ **rd**: raw biometric data (a image, a fresh capture)
- ▶ **bs**: biometric template to be compared with the reference
- ▶ **thr**: threshold according to a distance between templates
- ▶ **dec**: result of the matching (*decision*)

Analysis of Biometric Systems

An example of a biometric system

► the *Match-On-Card* technology



- storage of a reference inside a *secure element*
- the card performs the comparison
- the reference does not leave the card
- the terminal trusts the card

Analysis of Biometric Systems

Motivations

- ▶ Precedent work (BCLL'15; AL'14)
 - ▶ description of biometric systems within a formal model
 - ▶ formal reasoning about privacy properties
 - ▶ intuition: the architecture level is the right level to reason about privacy properties
 - ▶ assumption: the building blocks do their jobs properly
- ▶ This work
 - ▶ the precedent model is static
 - ▶ no runtime leakage is taken into account
 - ▶ we need such an extension for biometric systems

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A Formal Model for Privacy By Design

Formal Model for Privacy-By-Design

Reasoning about architecture

- ▶ Analysis of a system
 - ▶ components, localisation of data, trust assumptions, *etc.*
- ▶ Architecture language
 - ▶ Description of a system: a set of *architectural primitives*
 - ▶ e.g.: $Has_i(X)$, $Receive_{i,j}(\{S\}, \{X\})$, $Compute_i(X = T)$, ...
 $S ::= Attest_i(\{Eq\})$, $Eq ::= Pred(T_1, \dots, T_n)$
 - ▶ a description specifies each component, communication, computation, *etc.*
- ▶ Architecture semantics
 - ▶ semantics based on *traces* (aka sequences of events)
 - ▶ events are instantiations of architectural primitives
 - ▶ *architecture semantics*: the set of states reachable by the (compatible) traces

- ▶ A dedicated epistemic logic for *privacy properties*
 - ▶ following the “*deductive algorithmic knowledge*” paradigm
 - ▶ properties: *confidentiality* of data and *integrity* of computations
 - ▶ *semantics of a property P*: set of architectures satisfying *P*
- ▶ Axiomatics
 - ▶ aka a set of *deductive rules*, sound and complete with respect to the semantics e.g.:

$$\mathbf{H2} \frac{\text{Receive}_{i,j}(S, E) \in A \quad X \in E}{A \vdash \text{Has}_i^{\text{all}}(X)}$$

$$\mathbf{K1} \frac{\text{Compute}_i(X = T) \in A}{A \vdash K_i(X = T)}$$

Formal Model for Privacy-By-Design

Extension of the formal model

- ▶ Extending the architecture language
 - ▶ for each primitive, introduction of a **bound** on the number of instantiations
 - ▶ $Receive_{i,j}^{(n)}(\{S\}, \{X\})$, $Compute_i^{(n)}(X = T)$, ...
- ▶ Extending the traces of events
 - ▶ introduction of several sessions in the traces (a **Session** event)
 - ▶ enable the modelling of several successive sessions
 - ▶ enable the use of different values of the same variable
 - ▶ motivation: the information leakage due to the access of several values of the same variable across different sessions
 - ▶ introduction of a system re-initialization (a **Reset** event)

Formal Model for Privacy-By-Design

Extension of the formal model

- ▶ Adapting the trace semantics
 - ▶ the architecture semantics is always the set of reachable component states
 - ▶ the definition of the component states is adapted to the new events
- ▶ Extending the privacy logic
 - ▶ enable the reasoning about several access to the same variable
 - ▶ axiomatics for this extended logic:

$$\text{H2} \frac{\text{Receive}_{i,j}^{(n)}(S, E) \in A \quad X \in E}{A \vdash \text{Has}_i(X^{(n)})}$$

$$\text{H5} \frac{\text{Dep}_i(Y, \mathcal{X}) \quad \forall X^{(n)} \in \mathcal{X}: A \vdash \text{Has}_i(X^{(n)})}{A \vdash \text{Has}_i(Y^{(1)})}$$

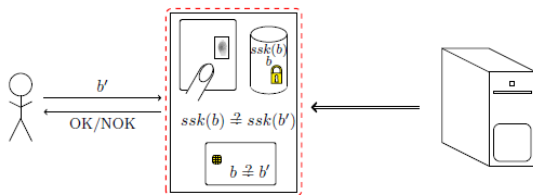
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Back to Biometric Systems

Analysis of Biometric Systems

Another biometric system

- ▶ Extending *Match-On-Card* to biometric identification (BCKK'09)



- ▶ terminal with extended functionalities
 - ▶ sensor + database + *smart card inside the terminal*
- ▶ biometric data are stored in two ways
 - ▶ a *secure sketch* for filtering; a *ciphertext* for identifying
- ▶ templates comparison carried out inside the card

Analysis of Biometric Systems

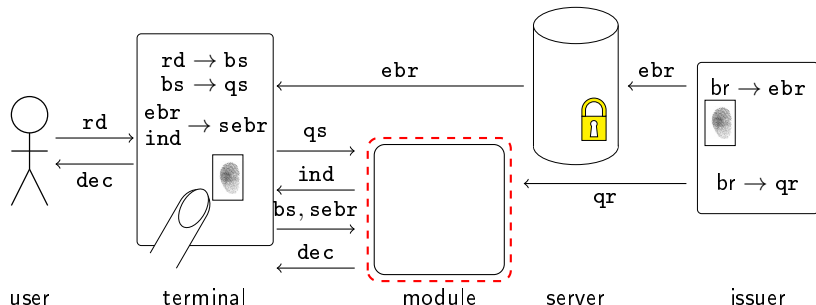
Inherent leakage in the black-box model

- ▶ information leakage
 - ▶ the result of the authentication *leaks information* about the quantizations ...
 - ▶ ... even if the functionality is seen as *black-box* (BCS'10)
- ▶ idea of the attack
 1. successive queries to the card with random values
 2. analysis of the indices queried by the card

Analysis of Biometric Systems

Application of the extended formal model

- ▶ description of this system and some variants inside the extended formal model
 - ▶ $Compute_T^{(n)}(sebr = EGet(ebr, ind)), \dots$



Analysis of Biometric Systems

Application of the extended formal model

- ▶ description variants with counter-measures
 - ▶ variant 1: ask all indices (no more dependencies on indices)
 - ▶ variant 2: block the number of queries before a critical bound
 - ▶ variant 3: re-initialization (re-encryption the database, permutation of the indices, etc.)
- ▶ then, analysis of the confidentiality of the secure sketches; e.g.:

$$\begin{array}{l} \nexists X : Dep_T(qr, X) \in A^{mi-e1} \quad \nexists j : Receive_T^{(n)}(S, \{qr\}) \in A^{mi-e1} \\ Has_T^{(n)}(qr) \notin A^{mi-e1} \quad \nexists T : Compute_T^{(n)}(qr = T) \in A^{mi-e1} \\ \hline \text{HN} \frac{\forall n : A \not\vdash Has_T(qr^{(n)})}{A \vdash Has_T^{none}(qr)} \end{array}$$

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Conclusion

- ▶ precedent work
 - ▶ description of biometric system architectures within a formal model
 - ▶ reasoning about privacy properties of the architectures
- ▶ our contribution
 - ▶ extending the formal model (semantics and epistemic logic) to catch the runtime leakage
 - ▶ important for biometric comparison
 - ▶ analysis of a variants a biometric system within this extension

thank you for your attention